### Hash Tables





Data Structures and Other Objects

- We discussed several ways of storing information in an array, and later searching for the information.
- Hash tables are a common approach to the storing/searching problem.
- This presentation introduces hash tables.

- The simplest kind of hash table is an array of records.
- This example has 701 records.



An array of records



#### 5033755339

- Each record has a special field, called its <u>key</u>.
- In this example, the key is an integer.

## [0] [1] [2] [3]



#### [4]

#### 5033755339

The number might be a person's campus phone number, and the rest of the record has information about the person.



 When a hash table is in use, some spots contain valid records, and other spots are "empty".

#### [0] [1] [2] [3] [4] [5]





In order to insert a new record, the <u>key</u> must somehow be <u>converted to</u> an array <u>index</u>.









Typical way create a hash value:

(Number mod 701)

What is (580625685 mod 701) ?





Hash with modular arithmetic:

(Phone# mod 701)

What is (5033706454 mod 701)?





The hash value is used for the location of the new record.









. . .

[0] [1] [2]

5033706862

The hash value is used for the location of the new record.



Here is another new record to insert, with a hash value of 2.



My hash value is [2].



 This is called a <u>collision</u>, because there is already another valid record at [2].

5033755314

5033706454

When a collision occurs, move forward until you find an empty spot.

 This is called a <u>collision</u>, because there is already another valid record at [2].

5033755314

5033706454

When a collision occurs, move forward until you find an empty spot.







This is called a <u>collision</u>, because there is already another valid record at [2].

5033755314

5033706454

伯

When a collision occurs, move forward until you find an empty spot.



 This is called a <u>collision</u>, because there is already another valid record at [2].

## The new record goes in the empty spot.



The data that's attached to a key can be found fairly quickly.



- Calculate the hash value.
- Check that location of the array for the key.

Not me.
(a)
(b)

[0]
[1]
[2]
[3]
[4]
[5]
[700]

[0]
[1]
[2]
[3]
[4]
[5]
[700]

[1]
[2]
[3]
[4]
[5]
[700]

[2]
[3]
[4]
[5]
[700]

[3]
[3]
[4]
[5]
[50370643]

[3]
[3]
[4]
[5]
[50370643]

[3]
[3]
[4]
[5]
[50370643]

[3]
[3]
[3]
[3]
[3]
[3]

[3]
[4]
[5]
[50370643]
[50370643]
[50370643]

[3]
[3]
[3]
[3]
[3]
[3]
[3]
[3]

[3]
[3]
[4]
[5]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]
[50370643]<

My hash



 Keep moving forward until you find the key, or you reach an empty spot. 5033706453

My hash value is [2].



 Keep moving forward until you find the key, or you reach an empty spot. 5033706453

Yes!

My hash value is [2].



5033706862

When the item is found, the information can be copied to the necessary location.

5033755314

5033706454

My hash value is [2].

5033706453

Yes!



#### Deleting a Record

Records may also be deleted from a hash table.



#### Deleting a Record

- Records may also be deleted from a hash table.
- But the location must not be left as an ordinary "empty spot" since that could interfere with searches.



#### Deleting a Record

- Records may also be deleted from a hash table.
- But the location must not be left as an ordinary "empty spot" since that could interfere with searches.
- The location must be marked in some special way so that a search can tell that the spot used to have something in it.



# Summary

- Hash tables store a collection of records with keys.
- The location of a record depends on the hash value of the record's key.
- When a collision occurs, the next available location is used.
- Searching for a particular key is generally quick.
- When an item is deleted, the location must be marked in a special way, so that the searches know that the spot used to be used.

Presentation copyright 2010 Addison Wesley Longman, For use with *Data Structures and Other Objects Using* C++ by Michael Main and Walter Savitch.

Some artwork in the presentation is used with permission from Presentation Task Force (copyright New Vision Technologies Inc) and Corel Gallery Clipart Catalog (copyright Corel Corporation, 3G Graphics Inc, Archive Arts, Cartesia Software, Image Club Graphics Inc, One Mile Up Inc, TechPool Studios, Totem Graphics Inc).

Students and instructors who use *Data Structures and Other Objects Using* C++ are welcome to use this presentation however they see fit, so long as this copyright notice remains intact.

